Let we have any set G (not necessary finite) consisting of the elements of any nature, i.e.  $G = \{a, b, c, ..., z, ...\}$ .

- 1. **Definition**. A set **G** is an algebraic <u>group</u> if it is equipped with a <u>binary operation</u> that satisfies four axioms:
- 1. Operation is closed in the set; for all a, b, there exists unique c in G such that a b = c.
- 2. Operation  $\bullet$  is associative; for all a, b, c in G:  $(a \bullet b) \bullet c = a \bullet (b \bullet c)$ .
- 3. Group G has an neutral element abstractly we denote by e such that  $e = e \cdot a = a$ .
- 4. Any element a in G has its inverse  $a^{-1}$  with respect to  $\bullet$  operation such that  $a \bullet a^{-1} = a^{-1} \bullet a = e$ .

For curiosity, can be said that group axioms seems very simple but groups and their mappings describes a very deep and fundamental phenomena in physics and other sciences. Among these mappings a special importance have mappings preserving operations from one group to another called isomorphisms, or homomorphisms and morphisms in general. Isomorphisms have a great importance in cryptography to realize a secure confidential *cloud computing*. It is named as *computation with encrypted data*. The systems having a homomorphic property are named as *homomorphic cryptographic systems*. They are under the development and are very useful in creation of secure e-voting systems, confidential transactions in blockchain and etc. We do not present there the construction of these systems and postpone it to the further issues of BOCTII, say in BOCTII.2. There we present one very important isomorphism example later when consider so called discrete exponent function (DEF).

T1. Theorem. If P is prime, then  $\mathcal{L}_p^* = \{1, 2, 3, ..., p-1\}$  where operation is multiplication mod p is a multiplicative group. Example:  $P = 11 \implies \mathcal{I}_p^* = \{1, 2, 3, ..., 10\}$ 

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Multiplication Tab. Z <sub>11</sub> *											$2.6 = 12 \mod 11 = 1$
*	1	2	(3)	4	5	6	7	8	9	10	_12 [11]
1	_(1	) 2	3	4	5	6	7	8	9	10	11 1
(2)	2	4	6	8	10	(1	3	5	7	9	1
3	3	6	9	1	) 4	7	10	2	5	8	
4	4	8	(1)	) 5	9	2	6	10	3	7	4.3 mad 11 = 12 mad 11 = 1
5	5	10	4	9	3	8	2	7	1	6	4.4° mad1 = (4/4) = 1)
6	6	1	7	2	8	3	9	4	10	5	# (1.1)
7	7	3	10	6	2	9	5	1	8	4	$4^{-1} = 3 \mod 11$
8	8	5	2	10	7	4	1	9	6	3	1 - 3 MIGALLY
9	9	7	5	3	1	10	8	6	4	2	5.9=45 mad 11=1=
10	10	9	8	7	6	5	4	3	2	1	45 111
											12

>> mod(4\*3,11) ans = 1

>> mulinv(4,11)

 $7/4 \mod 11 = 7*3 \mod 11 = 10$ 

ans = 3

Power Tab. Z <sub>11</sub> *						)	$(\epsilon a)$	Z10			
^	0	1	2	3	4	5	6	7	8	9	10
1	1	1	1	1	1	1	1	1	1	1	1
2	1	2	4	8	5	10	9	7	3	6	1
3	1	3	9	5	4	1	3	9	5	4	1
4	1	4	5	9	3	1	4	5	9	3	1
5	1	5	3	4	9	1	5	3	4	9	1
6	1	6	3	7	9	10	5	8	4	2	1
7	1	7	5	2	3	10	4	6	9	8	1
8	1	8	9	6	4	10	3	2	5	7	1
9	1	9	4	3	5	1	9	4	3	5	1
10	1	10	1	10	1	10	1	10	1	10	1

$$J_{M}^{*} = \{1,2,3,...,09\}$$
 $J_{10} = \{0,1,2,3,4,5,6,7,8,9\}$ 
 $DEF: J_{10} \longrightarrow J_{M}^{*}$ 
 $DEF_{2}(X) = 2^{X} \mod M = Q \in J_{M}^{*}$ 
 $\Rightarrow p=11;$ 
 $\Rightarrow g=2;$ 

>> mod exp(g,10,11) $\rightarrow$  mod exp(g,0,11) ans = 1ans = 1>> mod\_exp(g,1,11) >> mod exp(g,11,11)ans = 2ans = 2>> mod exp(g,12,11)>> mod\_exp(g,2,11) ans = 4ans = 4>> mod\_exp(g,13,11) >> mod\_exp(g,3,11) ans = 8ans = 8>> mod exp(g,14,11)>> mod\_exp(g,4,11) ans = 5ans = 5

Card 
$$(\mathcal{I}_{10}) = |\mathcal{I}_{10}| = 10$$
  
Card  $(\mathcal{I}_{n}^{*}) = |\mathcal{I}_{n}^{*}| = 10$   $\Rightarrow$  card  $(\mathcal{I}_{n}) = card (\mathcal{I}_{n}^{*})$ 

It is proved that: if p is prime, then there exists such numbers q that DEFg(X) provides 1- to-1 or bijective mapping.

Z & J.\* T2. Fermat (little)Theorem. If p is prime, then [Sakalauskas, at al.]

 $Z^{p-1} = Z^{\circ} = 1 \mod p$   $Z \mod p = Z \mod p$  $z^{p-1} = 1 \mod p$ 

How to find inverse element to z mod n?

Inverse elements in the Group of integers < Zp\*, • mod p> can be found using either

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How to find inverse element to z mod n?
>> mulinv(z.n)
Inverse elements in the Group of integers \langle \mathbf{Z}_{p}^{*}, \frac{\bullet_{\text{mod }p}}{\bullet_{\text{mod }p}} \rangle can be found using either
Extended Euclidean algorithm or Fermat theorem, or ...
Let we have z in Z_p^*, then to find z^{-1} mod p it can be done by Octave:
>> z m1=mulinv(z,p)
Let p is prime.
                                                                                                 >> p=genstrongprime(28)
Then p is strong prime if p=2q+1 where q=(p-1)/2 is prime as well.
                                                                                                 p = 144658379
                                                                                                 >> isprime(p)
Then g in \mathbb{Z}_{p}^{*} is a generator of \mathbb{Z}_{p}^{*} if and only if
                                                                                                 ans = 1
(iff) g^2 \neq 1 \mod p and g^q \neq 1 \mod p.
                                                                                                 >> q=(p-1)/2
                                                                                                 q = 72329189
For example, let p is strong prime and p=11, then one of the generators is g=2.
                                                                                                 >> isprime(q)
Verification method: g^2 \neq 1 \mod p and g^q \neq 1 \mod p.
                                                                                                 ans = 1
The main function used in cryptography is Discrete Exponent Function - DEF:
                                                                                                 >> g=2;
DEF_g(x) = g^x \mod p = a.
                                                                                                 >> mod_exp(g,2,p)
                                                                                                 ans = 4
                                                                                                 >> mod exp(g,q,p)
                                                                                                 ans = 144658378
Due to Fermat (little) theorem operations in exponents are performed mod (p-1).
g<sup>(a+b) mod (p-1)</sup> mod p
>> pp=11
pp = 11
>> gg=2
gg = 2
>> a=5
a = 5
>> b=9
b = 9
>> apb=mod(a+b,10)
apb = 4
>> g_apb=mod_exp(g,apb,11)
g_apb = 5
>> g_apb=mod_exp(g,14,11)
gapb = 5
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